

# Typing network communication

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Ivan Nikitin

University of Glasgow

# Types

```
/**  
 * Undocumented function  
 * @return void  
 */  
public static function test2() : void
```

Expected type 'float'. Found 'int'. intelephense(10006)

```
{  
    $sum = Prize::test3(1, 1);  
    echo $sum;  
}
```

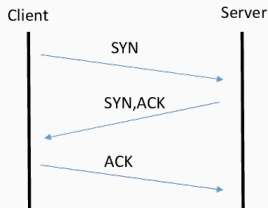
```
/**  
 * Undocumented function  
 * @param int $a  
 * @param int $b  
 * @return int  
 */  
private static function test3(int $a, float $b) : float  
{  
    return $a + $b;  
}
```

We can use static typing to:

- Protect from run-time errors.
- Document intention.
- Provide clear errors.

# Behavioural Types - Session Types

- Behaviour - the communication that occurs in terms of sequences of actions.
- TCP Handshake - Send a SYN, get a SYN-ACK, send an ACK.
- We can express this communication as a *type*.



# Behavioural Types - Session Types

$$\Gamma = \begin{array}{l} s[\text{server}] : \text{client} \oplus \text{syn}(\text{SYN}).\text{client}\&\text{synack}(\text{SYN-ACK}) \\ \quad \quad \quad .\text{client} \oplus \text{ack}(\text{ACK}).\text{end}, \\ s[\text{client}] : \text{server} \&\text{syn}(\text{SYN}).\text{server} \oplus \text{synack}(\text{SYN-ACK}) \\ \quad \quad \quad .\text{server}\&\text{ack}(\text{ACK}).\text{end} \end{array}$$

- $\oplus$  - Send (Select),  $\&$  - Receive (Offer),  
  . - continue as, **end** - end of communication.
- Each participant's local type is their view of the communication.
- E.g. from the server viewpoint - send a message of type **SYN** to the client, continue as receive a message of type **SYN-ACK** from the client, continue as send a message of type **ACK** to the client, continue as **end**.

# Behavioural Types - Session Types

An encoding of session types in a programming language gives us the ability to ensure that:

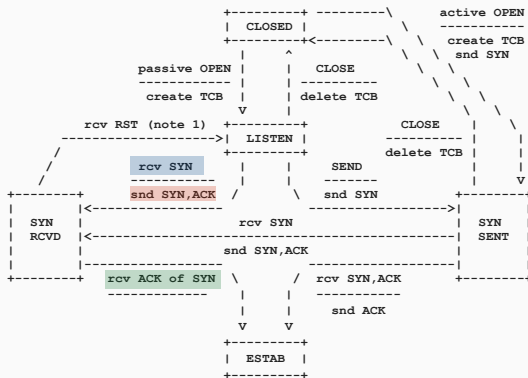
- Messages sent and received are in the expected order.
- Messages are of the expected type.
- Communication occurs between the expected participants.

# Behavioural Types - Session Types

An approach to encoding session types can be:

- Encode session type actions as a type. E.g.
  - $\oplus$  — *Select*  $\langle$  Message, Continuation  $\rangle$ ,
  - $\&$  — *Offer*  $\langle$  Message, Continuation  $\rangle$
- Channel parametrised over roles.  
*Channel*  $\langle$  ParticipantOne, ParticipantTwo  $\rangle$
- Channel has functions that use the session type.  
*select*(message, session\_type) —  $\rangle$  *SessionType*
- Continuation passing style ensures adherence to the sequence of actions. Parametrisation over roles ensures the channel is used correctly.

# Behavioural Types - Session Types



```
Γ = s[ss]:cs&syn(SynSet).
    cs @ syn_ack(SynAckSet).
    cs&(AckSet).
    cs @ fin(Fin),end
```

```
// Instantiate the session type of the TCP server.
type ServerSessionType = OfferOne<RoleClient, Syn,
    SelectOne<RoleClient, SynAck,
    OfferOne<RoleClient, Ack,
    SelectOne<RoleClient, FinAck,
    End>>>;

let st_server = ServerSystemSessionType::new();
// Create the session typed NetChannel
let mut chan = NetChannel::<RoleServer,
    RoleClient>::new(iter, tx, remote_addr);
// Recieve a SYN packet.
let (syn, cont) = chan.offer_one(st_server);
// Send the message along the channel, following our
// session type.
let cont = chan.select_one(cont, SynAck {...});
// Recieve a message of type ACK.
let (ack, cont) = chan.offer_one(cont);
// Send the FIN-ACK along the channel.
let cont = chan.select_one(cont, FinAck {...});
// Close
chan.close(cont);
```



# To Summarise

- We can encode the *behaviour* of a protocol as a *type*.
- Session types are a discipline used to express concurrent communication.
- An encoding of session types in a programming language leverages the compiler to make sure the protocol is implemented according to its session type.